

纯粹数学前沿

数论: 几何, 分析, 代数相交的地方

Introduction

(intro. 中文
后面 英文书)

数学中心问题:

解方程!

来源于物理: ordinary / partial differential equation.

例: Fluid mechanics: Navier-Stokes 方程.

Clay 数学所提出的千禧年大难题之一. (Millennium Pbs)

另一类重要的方程: 来源于“数”本身的方程.

实数 \rightarrow 自然数 \mathbb{N} .

$x+1=0$ 求解 \rightarrow 加入负数 $\rightarrow \mathbb{Z}$.

$3x=1$ 求解 \rightarrow 加入有理数 $\rightarrow \mathbb{Q}$.

$x^2=2$ 求解 \rightarrow 加入无理数 $\rightarrow \mathbb{R}$.

$x^2=-1$ 求解 \rightarrow 加入虚数 $\rightarrow \mathbb{C}$.



\rightarrow 多项式方程, 一元 n 次方程

Thm 所有多项式方程均在 \mathbb{C} 上有解.

仍然要问: 多项式方程 (或更一般的方程) 什么时候在 \mathbb{R} 中有解?
(实数)

数学分析 \rightarrow

① 有理数多项式方程什么时候在 \mathbb{Q} 中有解?

② 整数 \mathbb{Z} ?

② 称为丢番图 (Diophantine equation) 方程可解性问题.

Hilbert 第十问题: 能用一种由有限步构成的 ^{algorithm} 一般算法判断一个丢番图方程是否有解?

1970年, Yuri Matiyasevic 前苏联: 不存在这样的算法!

因此②很难.. 今天我们看①. (也很难)

有理系数多项式方程(组)什么时候在 \mathbb{Q} 中有解?

牵涉的方向: 数论: \mathbb{Z}, \mathbb{Q} , 数域 (\mathbb{Q} 的有限扩张) 代数几何: 多项式方程所定义的几何对象 (例: $x^2+y^2=1$ 圆) 算术代数几何

一些著名的例子:

"几何决定算术"

$\mathbb{P}_{\mathbb{C}}^2$ $(x:y:z)$ 复射影平面

给定一个齐次多项式 $P(x,y,z)$ degree = d .

$C = \{ (x:y:z) \in \mathbb{P}_{\mathbb{C}}^2 \mid P(x,y,z) = 0 \}$ 射影曲线 (复维数 1) Riemann surface 黎曼面 (实维数 2)

亏格 是它的一个几何量.

genus

" $\subset \mathbb{Q}$ "

Formula If C is smooth then $g(C) = \frac{1}{2}(d-1)(d-2)$ genus

From now on assume that C is smooth.

Rk. (1) 这可以作为光滑曲线亏格的定义.

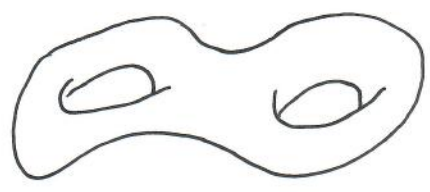
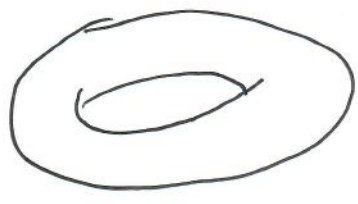
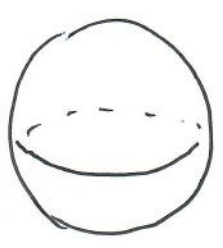
(2) 但这不是一个好的定义. 亏格作为一个"几何量"应该是 intrinsic 的.

同构的曲线应该有一样的亏格. 但 "degree" = 方程的次数. 同构的曲线对应的方程并不一样. ~~然而~~ 没有理由认为次数相同.

(3) “内蕴” intrinsic 的定义左是 C 的某个上同调群的维数。
(从几何体本身出发去定义, 而不是从方程出发去定义)

~~dim~~ $g(C) = \dim_C H^1(X, \mathcal{O}_X)$

几何直观: 黎曼面 / (代数) 射影曲线



亏格 $g = 0$

1

2

算术性质: 当定义 C 的齐次式 $P(x, y, z)$ 的系数在 \mathbb{Q} 中时,
齐次式 P 是否有有理数解?

~~例 Fermat 大定理 (Wiles) 1993. $P(x, y, z) = x^n + y^n - z^n$ 没有非平凡解!~~

以下为

几何决定算术的三个定理:

神奇之处在于 输入 是一个几何量, g . 亏格. 是 P 的所有复数解 (组成一个复流形——黎曼面) 的一个纯几何不变量.

而 输出 的结论是关于 P 在 \mathbb{Q} 中解方程的可能性!

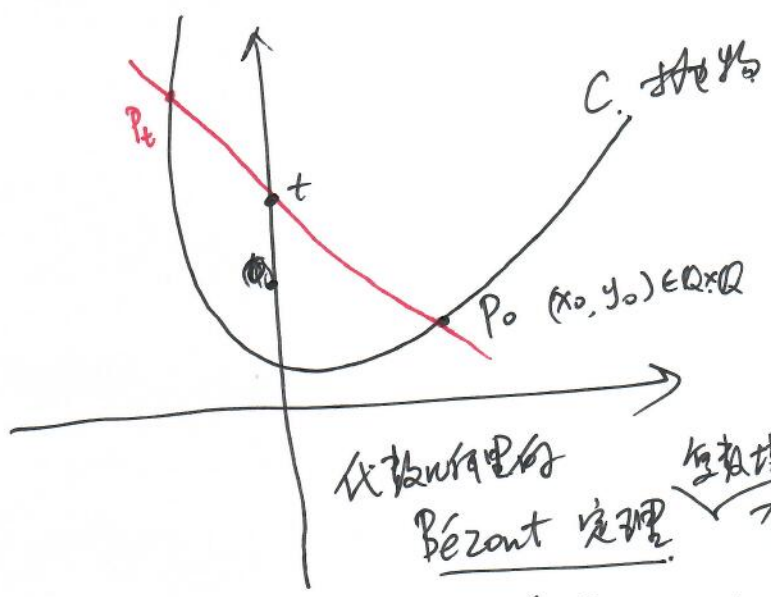
~~Thm 1. $g(C) = 0$ 而且 C 有一个有理点 (即 P 有一个有理数解)~~

那么

Thm 1. $g(C) = 0$ and if C has one rational point
(i.e. P has a solution in \mathbb{Q})

Then C has as many as rational points as \mathbb{P}^1
Rk. In particular, the number of rational points is infinite.

Proof. $g = 0 \Leftrightarrow d = 2$. 二次曲线/圆锥曲线 $\left\{ \begin{array}{l} \text{抛物线} \\ \text{椭圆} \\ \text{双曲线} \end{array} \right.$



Fix a rational point P_0 of C .

任取 y 轴上一个有理数 t .
连直线 tP_0 .

代数几何的 Bézout 定理 $\left\{ \begin{array}{l} \text{有理数域上} \\ \text{直线} \end{array} \right.$ (一次曲线) 与二次曲线在射影平面内相交 ~~必有~~ 的交点数为 $2 \times 1 = 2$.

C 与 --- 相交得 P_t 与 P_0

$t \in \mathbb{Q}$.
 $P_0(x_0, y_0) \in \mathbb{Q} \times \mathbb{Q}$ } \Rightarrow 直线斜率 $\in \mathbb{Q}$

\Rightarrow 代入方程 P (有理系数) P_t 的坐标 $\in \mathbb{Q}$.
即 P_t 是 C 的一个有理点

(y 轴) $\mathbb{P}^1 \rightarrow C$
 $t \mapsto P_t$ 是一个双射.

#.

(Mordell-Weil)

Thm 2. $g(C) = 1$, if C has one rational point.

Then the set of rational points of C is a finitely generated abelian group \mathcal{O} .

Rk. (1) Mordell is \mathbb{Q} \mathbb{Z}

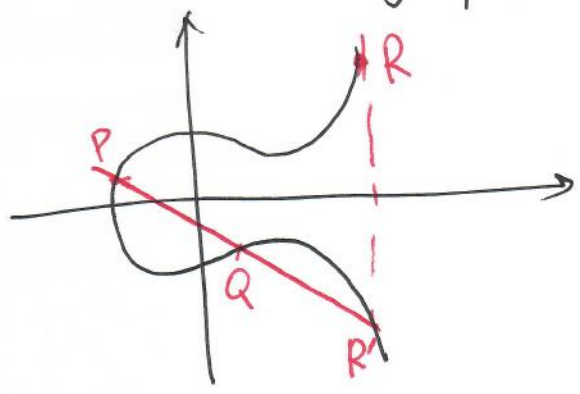
(2) Weil is \mathbb{Q} : finite extensions of \mathbb{Q} , e.g. $\mathbb{Q}(i)$, $\mathbb{Q}(\sqrt{2})$

(3) In this case, we say that C is an elliptic curve.
Weierstrass equation: C is isomorphic to $y^2 = x^3 + ax + b$ $\Delta = 4a^3 + 27b^2 \neq 0$
 $zy^2 = x^3 + ax^2 + bx^2$ $(g=1:0) \in \text{elliptic}$.

(4) the number of rational points can be finite or infinite.

~~Structure theorem of finitely generated abelian~~

(5) the set of rational points $\mathcal{O} \subset C(\mathbb{Q})$ (or $C(k)$)
 k number field
is an abelian group!



$y^2 = x^3 + ax + b$
 C degree 3 curve

Bezant \Rightarrow intersection = 3 points

$P+Q := R$. check that this is a group!

Commutative group.

$\mathcal{O} = (0:1:0)$
~~无意义~~

(6) Structure thm for finitely generated abelian group:

$C(\mathbb{Q}) \cong F \oplus \mathbb{Z}^r$
finitely gp.

BSD Conjecture: (Clay 23rd problem)
 $r = \text{ord}_{s=1} L(C, s)$

L function of the elliptic curve.

(7) The proof uses Fermat's descent method. 费马递降法.


and height function (高度) ~~for~~ to measure the complexity of a rational points.

Thm 3 (Faltings) $g(C) \geq 2$ Then $C(\mathbb{Q})$ (or $C(k)$) is finite.

C has at most finitely many rational points.


Fields medal 1986

The proof uses Faltings' height. \rightarrow 袁新喜 2019年5月在科大的短课程. (科大网络的课程有视频).

为什么
↓
算术
 $g=0$

 $|C(\mathbb{Q})| = \infty$

$g=1$

 $C(\mathbb{Q})$ finitely generated abelian group.

$g \geq 2$

 $|C(\mathbb{Q})| < \infty$

Fermat's last theorem (Wiles 1993)

$C: x^n + y^n = z^n \quad n \geq 3$ has no non-trivial solution.

$n=3, 4$ proved by Fermat.

$n=5. \quad g = \frac{1}{2}(n-1)(n-2) \geq 2. \quad$ Faltings \Rightarrow only finitely many solution!

Wiles \Rightarrow no non-trivial solution! much stronger!

proof uses $\left\{ \begin{array}{l} \text{elliptic curve} \\ \text{modular form (number theory)} \\ \text{galois representation.} \end{array} \right.$

Q32. How to study rational solutions of polynomials?

K = number field = finite extension of \mathbb{Q}

X_K = algebraic variety = a set of polynomials with coeff. in K

\downarrow field extension L/K $X(L) = \{ \}$ the set of solutions in L of the polynomials.
= set of rational points

Suppose that X is smooth (i.e. $X(\mathbb{C})$ is a smooth complex manifold)

example: X defined by $P(x,y) = x^2 + y^2 + 1$ over $K = \mathbb{Q}$.

~~$X(\mathbb{Q}) = \emptyset, X(\mathbb{R}) = \emptyset, X(\mathbb{C}) = \emptyset.$~~

$X(\mathbb{C}) \neq \emptyset$ \mathbb{C} = algebraically closed.

$X(\mathbb{Q}) = \emptyset$ why? Since $X(\mathbb{R}) = \emptyset.$

real analysis $\Rightarrow \begin{cases} x^2 \geq 0 \\ y^2 \geq 0 \end{cases} \Rightarrow X(\mathbb{R}) = \emptyset. \Rightarrow X(\mathbb{Q}) = \emptyset$
 $\boxed{\mathbb{Q} \subseteq \mathbb{R}}$

advantage of \mathbb{R} : ① can do real analysis (i.e. can take limit. complete)
Cauchy sequences are convergent

② \mathbb{R} is not far from \mathbb{Q} .

$\mathbb{Q} \subseteq \mathbb{R}$ is dense

\mathbb{R} solutions may be approximated by \mathbb{Q} solutions.

① + ②: \mathbb{R} is a completion of \mathbb{Q} .

Natural question: other completions?

\mathbb{Q} , $a, b \in \mathbb{Q}$. $d(a, b) = |a - b|_{\infty} := |a - b|$ ^{absolute value} ~~$d(a, b)$~~ is a distance on \mathbb{Q}

- ① $d(a, b) \geq 0 \quad \forall a, b$; $d(a, b) = 0$ iff $a = b$.
- ② $d(a, b) = d(b, a)$
- ③ $d(a, b) + d(b, c) \geq d(a, c)$

Completion.
 add limits of all Cauchy sequence $\rightarrow \mathbb{R}$

Define a new distance:

$p \in \mathbb{Z}$ prime number.

$\forall n \in \mathbb{Z}$. $v_p(n) := r$ if $p^r | n$ but $p^{r+1} \nmid n$.

$\forall \frac{m}{n} \in \mathbb{Q}$ $v_p(\frac{m}{n}) := v_p(m) - v_p(n)$ well-defined.

$$|\frac{m}{n}|_p := p^{-v_p(\frac{m}{n})}$$

$$d_p(\frac{m_1}{n_1}, \frac{m_2}{n_2}) := |\frac{m_1}{n_1} - \frac{m_2}{n_2}|_p$$

- ① ✓ ② ✓, ③': $d_p(a, c) \leq \max(d_p(a, b), d_p(b, c))$
- \Rightarrow ③

p -adic distance.
 $p \nmid$

example

$p = 3$

$n_1 = 36 = 3^2 \times 2^2$
 $n_2 = 3$
 $n_3 = 27 = 3^3$

$v_p(n_1) = 2$, $|n_1|_p = \frac{1}{9}$ ✓
 $v_p(n_2) = 1$, $|n_2|_p = \frac{1}{3}$ ✗
 $v_p(n_3) = 3$, $|n_3|_p = \frac{1}{27}$ ✗

$$\mathbb{Q} \xrightarrow[\text{I. } | \cdot |_p]{\text{Completion}} \mathbb{Q}_p$$

$\mathbb{Q} \subseteq \mathbb{Q}_p$ can do p-adic analysis on \mathbb{Q}_p .

Thm (Ostrowski 1916) Every non-trivial absolute value on \mathbb{Q} is equivalent to either $| \cdot |_p$ or $| \cdot |_\infty$.

→ We should consider \mathbb{R} and all \mathbb{Q}_p .

分析表述

$$\left. \begin{array}{l} \mathbb{Q} \xrightarrow{\text{Completion}} \mathbb{Q}_p = \{p\text{-adic numbers}\} \\ \mathbb{Z} \xrightarrow{\quad} \mathbb{Z}_p = \{p\text{-adic integers}\} \end{array} \right\}$$

代数表述:

$$\mathbb{Q} = \text{Frac}(\mathbb{Z}), \quad \mathbb{Q}_p = \text{Frac}(\mathbb{Z}_p)$$

fraction field

~~代数~~

$$\mathbb{Z}_p = \varprojlim_n \mathbb{Z}/p^n \mathbb{Z} \subseteq \prod_{n \geq 1} \mathbb{Z}/p^n \mathbb{Z} \quad \boxed{\mathbb{Q}_p = \mathbb{Z}_p \otimes_{\mathbb{Z}} \mathbb{Q} = \text{Frac}(\mathbb{Z}_p)}$$

$:= \{ (a_n)_{n \geq 1} \mid a_n \in \mathbb{Z}/p^n \mathbb{Z} \text{ s.t. } a_n \text{ mod } p^n = a_{n+1} \text{ mod } p^n \forall n \geq 1 \}$

Consider ~~\mathbb{Z}_p -solutions~~ solutions in \mathbb{Z}_p or in \mathbb{Q}_p .

→ \mathbb{Z}_p -solutions more or less

- ~~mod p~~ solutions
- mod p^2 solutions
- mod p^n solutions.

p-adic analysis:

Hensel's Lemma:

$$f(x) \in \mathbb{Z}[x], \quad k \in \mathbb{N}, \quad r \in \mathbb{Z} \text{ st. } f(r) \equiv 0 \pmod{p^k}$$

(i.e. $|f(r)|_p \leq \frac{1}{p^k}$)

$$m \in \mathbb{N}, \quad m \leq k$$

If $f'(r) \not\equiv 0 \pmod{p}$ (i.e. $|f'(r)|_p = 1$)

Then $\exists s \in \mathbb{Z}$ st. $\begin{cases} f(s) \equiv 0 \pmod{p^{k+m}} & \text{(i.e. } |f(s)|_p \leq \frac{1}{p^{k+m}} \text{)} \\ s \equiv r \pmod{p^k} & \text{(i.e. } |s-r|_p \leq \frac{1}{p^k} \text{)} \end{cases}$

Moreover, s is unique mod p^{k+m} .

f has a mod p^k solution $\xRightarrow{f \text{ is "good"}}$ f has a mod p^{k+m} solution.

\rightarrow get mod $p, \text{ mod } p^2, \dots, \text{ mod } p^n \dots$ solution

\rightarrow get \mathbb{Z}_p -solution.

Geometric version: (数论分析: 牛顿迭代, 收敛速度与精度) $\text{req: } \begin{cases} f=0 \\ f' \neq 0 \end{cases} \Rightarrow$ Jacobi-Newton 迭代法. 充分性.

If X is smooth mod p , then $X(\mathbb{F}_p) \neq \emptyset \Rightarrow X(\mathbb{Z}_p) \neq \emptyset$
 \rightarrow good reduction mod p . $X(\mathbb{Q}_p)$

Rk: (1) $X(\mathbb{F}_p)$ is "easy" to compute (by computer!)

(2) Hensel's lemma says: easy to get \mathbb{Q}_p -solution if ... and reduction mod p

(3) X has good reduction mod p for all but finitely many p .

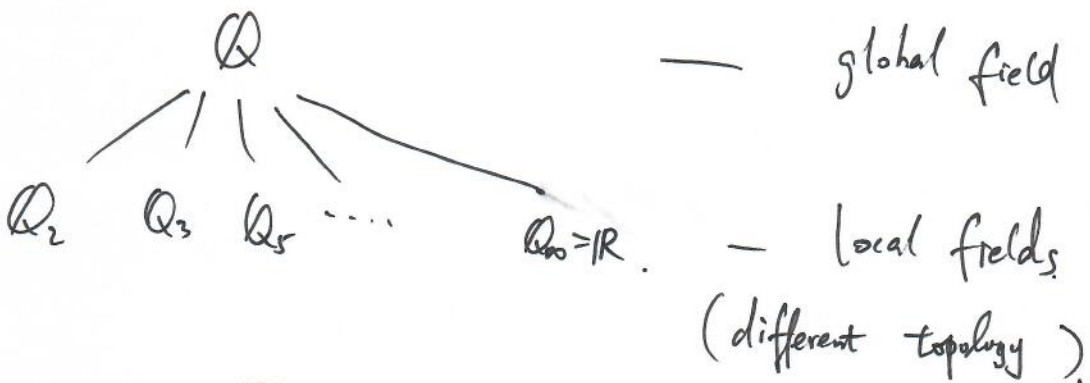
example

X defined by $P(x,y) = x^2 + 45y^2 - 75$

has good reduction mod p if $p \neq 3, 5$.

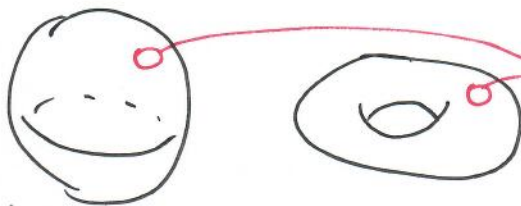
$p=3$ or 5 : ~~X mod p~~ X mod p is ~~not~~ defined by $x^2 = 0$ mod p .
double point, not smooth.

§3 global and local in number theory.



in geometry.

Why call it local



locally looks the same

globally different. $g=0, g=1$.

In alg. geo. $\text{Spec}(\mathbb{Z})$ global geom. object.

prime numbers are points on $\text{Spec}(\mathbb{Z})$.
each p is a "local object".

In number theory

$$\exists \mathbb{Q} \text{ solutions (global solutions)} \Rightarrow \exists \mathbb{Q}_p \text{ solutions (local solutions)}$$

Question

$\Omega = \{\text{primes}\} \cup \{\infty\}, \mathbb{Q}_\infty = \mathbb{R}$.

Def. Hasse principle holds if $X(\mathbb{Q}_p) \neq \emptyset \forall p \in \Omega \Rightarrow X(\mathbb{Q}) \neq \emptyset$.
(Local-global principle)

Thm (Hasse-Minkowski) If X is defined by quadratic form.

then ~~Hasse~~ local-global principle holds for X .

Rk. (1) Hasse proved it for \mathbb{Q} . Minkowski for a general number field.

(2) ~~An~~ elementary proof \rightarrow J.P. Serre ~~in~~ « A Course in Arithmetic ».

(3) We ~~are~~ are going to give a "proof" for \mathbb{Q} a special case.

X defined by

$$P(x,y,z) = x^2 + ay^2 + bz^2$$
 quadratic form. $a, b \in \mathbb{Q}^*$

Def (Hilbert Symbol) K/\mathbb{Q} field extension.

$$(a, b)_K := \begin{cases} 1 & \text{if } X(K) \neq \emptyset \\ -1 & \text{if } X(K) = \emptyset. \end{cases}$$

Notation : $(a, b)_p := (a, b)_{\mathbb{Q}_p}$
 $(a, b)_\infty := (a, b)_{\mathbb{R}}$
 $(a, b) := (a, b)_{\mathbb{Q}}$

Want to prove Hasse principle : $\left. \begin{matrix} (a, b)_p = 1 \ \forall p \\ (a, b)_\infty = 1 \end{matrix} \right\} \Rightarrow (a, b) = 1$

Before ~~to~~ prove, get a feeling of local and global in Number theory.

Question! for each $p \in \mathbb{R}$ ~~(and ∞)~~ fix $n_p \in \{\pm 1\}$.

Does there exist $a, b \in \mathbb{Q}^*$ st. $(a, b)_p = n_p \ \forall p$?
 $(a, b)_\infty = n_\infty$?

Necessary conditions: (a local condition and a global condition)

Thm. For any $a, b \in \mathbb{Q}^*$

(C1) $(a, b)_p = 1$ for almost all p .

(local property, follows from Hensel's lemma for good reduction primes)

(C2) product formula.

~~$\prod_{p \in \mathbb{Z}}$~~ $\prod_{p \in \mathbb{Z}} (a, b)_p = 1$

- ① " \prod " makes sense since (1)
- ② global property: the values $(a, b)_p$ are not independent they have at least this relation $\prod (a, b)_p = 1$
- ③ this follows from quadratic reciprocity law of Gauss $=$ 二次互反律

We reduce Question 1 to

Question 2: Are these conditions C1, C2 sufficient conditions for Question 1?

We are going to answer to Q2 and prove the Hasse principle for X .

Def. K field, a K -algebra A is a ring A containing K . $K \subseteq A$. (in particular, A has identity element $1_A = 1_K$)

A may be non-commutative.

~~From now on suppose that $\dim_K A < \infty$ (viewed as a K -vector space)~~

(i.e. K -algebra = ring + K -vector space structure. all operations are compatible)

Suppose that $\dim_K A < \infty$. from now on.

We say that A is a simple algebra if it has no non-trivial

A is a central simple algebra if $\text{center}(A) = K$.

example: $A = M_n(K)$ $n \times n$ matrices

Prop. A : K -algebra. TFAE.

(1) A is a central simple algebra

(2) $A \otimes_K K^s \cong_{K^s\text{-alg}} M_n(K^s)$ (K^s ~~is~~ separable closure)

(3) \exists finite extension L/K s.t. $A \otimes_K L \cong_{L\text{-alg}} M_n(L)$ (之后再证)
~~(4) (Wedderburn) $A \cong M_n(D)$ where D is a division algebra.~~ ← representation theory of algebras. 2/3/19

Rk: (1) \otimes = tensor product (homological algebra, commutative algebra)

Coefficient $\in K \rightsquigarrow$ in L (L/K)

e.g. $M_n(K) \otimes_K L = M_n(L)$ for matrix algebra.

(2) \otimes central simple algebra more or less matrix algebra.

↑
after a finite separable extension

It is a "twist" of the matrix algebra

example: \mathbb{H} : Hamilton's quaternions algebra over $K = \mathbb{R}$

As K -vector space $\mathbb{H} = 1 \cdot K \oplus i \cdot K \oplus j \cdot K \oplus k \cdot K$

basis $\{1, i, j, k\}$ $\omega_i \in \mathbb{R} K = \mathbb{R}$

product in \mathbb{H} is given by $i^2 = -1, j^2 = -1, ij = -ji = k$
 ~~$ij = k$~~

Prop \mathbb{H} is a division algebra (ie. non-zero elements are invertible)

proof: Norm map $N: \mathbb{H} \rightarrow \mathbb{R}$

$$q = x + yi + zj + tk, \quad x, y, z, t \in \mathbb{R}.$$

$$\begin{aligned} N(q) = q \cdot \bar{q} &= (x + yi + zj + tk)(x - yi - zj - tk) \\ &= \dots \\ &= x^2 + y^2 + z^2 + t^2 \in \mathbb{R} \end{aligned}$$

$$q \neq 0 \iff \text{one of } x, y, z, t \neq 0 \iff \begin{matrix} \uparrow \\ x, y, z, t \in \mathbb{R} \end{matrix} \iff \begin{matrix} \downarrow \\ N(q) \neq 0 \end{matrix}$$

$$\iff q^{-1} = \frac{\bar{q}}{N(q)} \quad \#.$$

Rk. The proof use the fact that ^{coefficients} $x, y, z, t \in \mathbb{R}$.

after tensor with \mathbb{C} . $\text{coeff} \in \mathbb{C}$, the same proof fails!

Indeed.

$\mathbb{H} \otimes_{\mathbb{R}} \mathbb{C}$	\cong	$M_2(\mathbb{C})$
1	→	$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$
i	→	$\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$ $\begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}$
j	→	$\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$
k	→	$\begin{pmatrix} 0 & -1 \\ -1 & 0 \end{pmatrix}$

(\Rightarrow $\mathbb{H} \otimes_{\mathbb{R}} \mathbb{C}$ is a twist of the matrix algebra.
i.e. a central simple algebra)

Rrk in $M_2(\mathbb{C})$ not all non-zero elements are invertible.

Norm map $\otimes_{\mathbb{R}} \mathbb{C} \rightsquigarrow \det$.

Generalization $a, b \in K^*$.

$Q_{a,b} := 1 \cdot K \oplus i \cdot K \oplus j \cdot K \oplus k \cdot K$ as vector space.

$$\begin{aligned}
 i^2 &= a \\
 j^2 &= b \\
 ij &= -ji = k.
 \end{aligned}$$

$Q_{a,b}$: quaternion algebra.

Thm (Wedderburn) (\leftarrow representation theory for finite dimensional semi-simple S algebras.)

$Q_{a,b}$ is either ① a division algebra or ② isomorphic to $M_2(K)$ (② definition split)

relation with solution over $K = \mathbb{Q}, \mathbb{Q}_p, \mathbb{R}, \dots$

Thm TFAE.

- (1) ~~$P(x,y,z)$~~ $P(x,y,z) = x^2 - ay^2 - bz^2$ has non-trivial solution in K
- (2) $(a,b)_K = 1$
- (3) $Q_{a,b}$ splits over K .

[key point: norm map ~~$N(x+yi+zj)$~~ $N(x+yi+zj) = x^2 - ay^2 - bz^2$]

Question 2 \iff " $Q_{a,b}$ splits over $\mathbb{Q}_p \forall p \in \mathbb{R}$ " \implies $Q_{a,b}$ splits over \mathbb{Q} ? "

Now we can use powerful tools ~~from~~ from algebra. ~~and number~~
(actually, from algebraic number theory)

Brauer group :

Def. $BrK = \{ \text{central simple algebra over } K \} / \sim$

\sim : equivalent relation. $A \sim B \stackrel{\text{def.}}{\iff} \exists m, n \in \mathbb{N}$ st. $M_n(A) \cong M_m(B)$ as K -algebrs.

(e.g. $A = M_r(K) \sim B = K$
take $n=1, m=r$.)

~~BrK~~ is an abelian group. (Brauer group of K)

product : $A \otimes_K B$

identity elemt : K , $A \otimes_K K \cong A$.

inverse : $A \otimes_K A^{op} \cong M_n(K) \sim K$.

(A^{op} = opposite ring of A
反环 i.e. $a \cdot b = ba$ in A
in A^{op})

Now quaternion algebras are central simple algebras,
they are elements in BrK .

$Q_{a,b} \in BrK$. \otimes ($Q_{a,b} \in {}_2 BrK$ 2-torsion ~~in~~ part.)

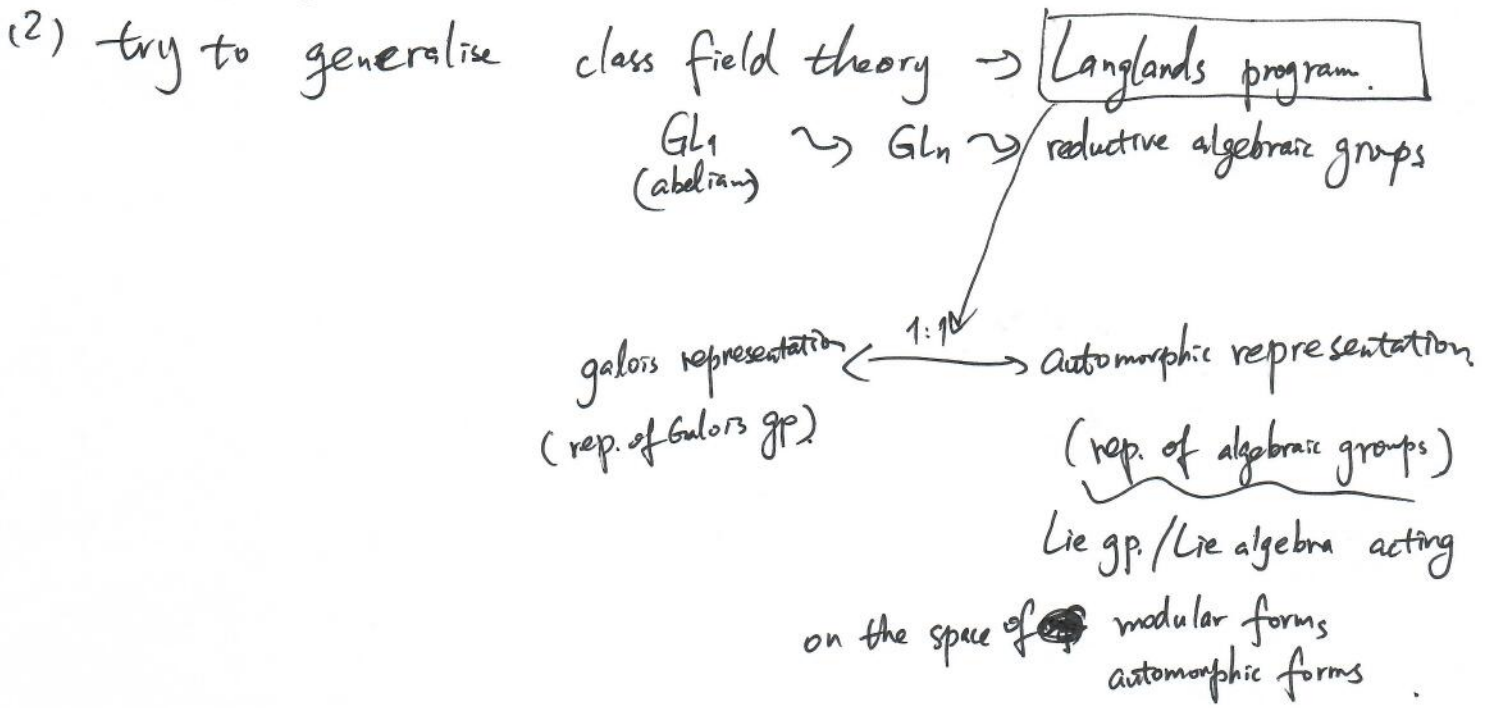
Thm (Global class field theory — algebraic number theory)

$$0 \rightarrow Br\mathbb{Q} \xrightarrow{\varphi} Br\mathbb{R} \oplus \left(\bigoplus_{\mathbb{P}} Br\mathbb{Q}_{\mathbb{P}} \right) \xrightarrow{\psi} \mathbb{Q}/\mathbb{Z} \rightarrow 0 \quad (*)$$

is an exact sequence.
正合列

Tate's thesis (harmonic analysis over local fields and number fields)
 homological method (cohomology of groups, Galois cohomology)

Rk. (1) class field theory (Takagi, E. Artin) is a generalization of Gauss' reciprocity law.



Back to This

$$\text{Br } \mathbb{R} \xrightarrow{\text{inv}_{\mathbb{R}}} \mathbb{Q}/\mathbb{Z}$$

$$\text{Br } \mathbb{Q}_p \xrightarrow{\text{inv}_p} \mathbb{Q}/\mathbb{Z}$$

$$\mathbb{Q}_{a,b} \longmapsto \begin{cases} 0 \pmod{\mathbb{Z}}, & \text{if } \mathbb{Q}_{a,b} \text{ splits in } \mathbb{Q}_p, (a,b)_p = 1 \\ \frac{1}{2} \pmod{\mathbb{Z}}, & \text{otherwise, } (a,b)_p = -1 \end{cases}$$

product formula $\prod_{p \in \mathbb{P}} (a,b)_p = 1 \Rightarrow$ ψ is a complex, i.e. $\psi \circ \varphi = 0$ (for the 2-torsion part)

~~exactness~~ exactness at the middle: $\ker(\psi) = \text{im}(\varphi)$ means:

$n_p \in \mathbb{Z} \setminus \{1\}$, \mathbb{Q} almost all 0, and $\prod_{p \text{ prime}} n_p = 1$

then $(n_p)_{p \in \mathbb{P}} \in \ker(\psi)$

$\Rightarrow \exists (a,b) \in \text{Br } K$ st. $(a,b)_p = n_p \quad \forall p$ prime or ∞ .

This answers to Question 2.

exactness ~~at~~ on the left means φ is injective:

$(a,b)_p = 1 \quad \forall p \in \mathbb{Z} \implies (a,b) = 1$
 ~~$(a,b)_p = 1$~~

i.e. local-global principle holds for X (defined by ~~$P(x,y,z) = x^2 + y^2 + z^2$~~
 $P(x,y,z) = x^2 - ay^2 - bz^2 = 0$)

§4. failure of Hasse principle.

quadratic form \checkmark .

~~qu~~ cubic form X .

Selmer: X defined by $P(x,y,z) = \del{3x^3 + 4y^3 + 5z^3} = 0$
 $3x^3 + 4y^3 + 5z^3$

has solutions in all \mathbb{Q}_p and \mathbb{R} .

but no solution in \mathbb{Q} .

Another easy counter example:
 X :

$P(x) = (x^2 - 13)(x^2 - 17)(x^2 - 13 \cdot 17)$

13, 17, 13x17 are not squares $\implies X(\mathbb{Q}) = \emptyset$.

$X(\mathbb{R}) \neq \emptyset \quad x = \sqrt{13} \checkmark$.

$2^2 = 17 \pmod{13} \implies X(\mathbb{F}_{13}) \neq \emptyset \xrightarrow{\text{Hensel}} X(\mathbb{Q}_{13}) \neq \emptyset$

$8^2 = 13 \pmod{17} \implies X(\mathbb{F}_{17}) \neq \emptyset \xrightarrow{\text{Hensel}} X(\mathbb{Q}_{17}) \neq \emptyset$

for $p \neq 13, 17$. Legendre symbol $\left(\frac{13}{p}\right) \cdot \left(\frac{17}{p}\right) = \left(\frac{13 \cdot 17}{p}\right)$

$\implies \forall (\mathbb{F}_p) \neq \emptyset$

~~can not~~ one of these must be 1.

§5 Weak approximation.

existence of solution, \rightsquigarrow how many solution.

$$X(\mathbb{Q}) \subset \prod_{p \in \mathbb{R}} X(\mathbb{Q}_p)$$

product topology

$p = \infty$
 $\mathbb{Q}_\infty = \mathbb{R}$

Weak approximation for X : $X(\mathbb{Q})$ dense in $\prod X(\mathbb{Q}_p)$.

means many \mathbb{Q} -solutions.

example. $X = P^1$

weak approx \Leftrightarrow chinese remainder theorem.

it may fail: $X:$

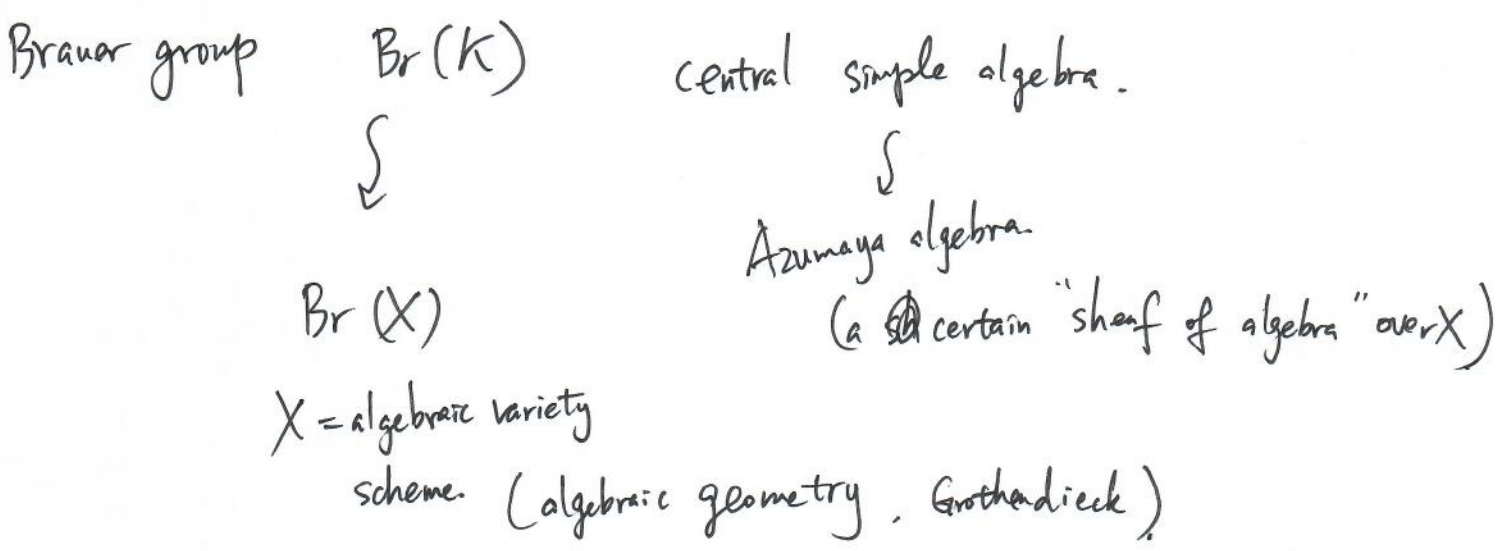
$$x^2 - 2y^2 = -(z^2 + 3)(z^2 - 3)$$
~~$$x^2 - 2y^2 = (z^2 - 3)(z^2 + 3) \quad / \mathbb{Q}.$$~~

$$\phi \neq X(\mathbb{Q}) \subsetneq \prod_{p \in \mathbb{R}} X(\mathbb{Q}_p) \quad (\mathbb{Q}_3 \text{ St-Br obs})$$

$(x, y, z) = (3, 0, 0)$

§6. Brauer-Manin obstruction

Aim: to explain the failure of Hasse principle and Weak approximation.



the Azumaya alg. definition is not a good definition.

Grothendieck developed étale cohomology

$$Br(X) := H_{\text{ét}}^2(X, G_m)$$

\uparrow use étale topology on X. \uparrow sheaf on X

This is good: functorial. 例子: 范畴语言 (Category language 范畴论的抽象语言)

i.e. $X \xrightarrow{f} Y$ map (morphism) induces $f^*: Br(Y) \rightarrow Br(X)$

algebraic geometry language: $X(k) = \text{Hom}(\text{Spec } k, X)$

rational points are maps between certain geometric objects!

$x \in X(k)$

$x: \text{Spec } k \rightarrow X$

induces $x^*: Br X \rightarrow Br k$
 $b \mapsto x^*(b) =: b(x)$

Yu. I. Manin (1970 ICM 国际数学家大会)

$$Br X \times \prod_{p \in \mathbb{Z}} X(\mathbb{Q}_p) \rightarrow \mathbb{Q}/\mathbb{Z}$$

$$b, (x_p)_{p \in \mathbb{Z}} \mapsto \sum_p \text{inv}_p(b(x_p))$$

$$0 \rightarrow Br \mathbb{Q} \rightarrow \bigoplus_{p \in \mathbb{Z}} Br(\mathbb{Q}_p) \rightarrow \mathbb{Q}/\mathbb{Z} \rightarrow 0$$

global local

→ global theorem says $0 \rightarrow \mathbb{Q}/\mathbb{Z} \rightarrow \dots$

$$X(\mathbb{Q}) \subset \overline{X(\mathbb{Q})} \subset \left[\prod_p X(\mathbb{Q}_p) \right]^{\text{Br}} \subseteq \prod_p X(\mathbb{Q}_p)$$

\swarrow closed.
 Br
 \parallel
 $\{ (x_p) \mid (x_p) \perp b \ \forall b \in \text{Br } X \}$

Rk (1) $\neq \Rightarrow$ weak approximation fails!

2) $\left[\prod_p X(\mathbb{Q}_p) \right]^{\text{Br}} = \emptyset \Rightarrow X(\mathbb{Q}) = \emptyset$ even if $\prod_p X(\mathbb{Q}_p) \neq \emptyset$.
 \Rightarrow local-global principle fails!

This is called the Brauer-Manin obstruction

This explains ~~many~~ failure of local-global principle / weak approx.
 for many algebraic ~~variety~~ varieties.

example. elli genus 1 curves. (e.g. $C: 3x^3 + 4y^3 + 5z^3 = 0$)

$E = \text{Jac}(C)$ Jacobian variety of C .
 \hookrightarrow elliptic curve.
 obstruction lies in $\text{III}^1(E^\vee) \in \text{Br}(E^\vee)$ ($E^\vee = \text{dual of } E$)
 Tate-Shafarevich group.

Conjecture (Colliot-Thélène et al.)

For rationaly connected varieties, the Brauer-Manin obstruction controls the failure of Hasse principle and weak approximation.

(1) Rationally connected: ^{有理连通.} geometric condition.:

$X(\mathbb{C})$: complex manifold. every 2 points can be connected by a projective line.

$\forall P_1, P_2 \in X(\mathbb{C}), \exists f: \mathbb{P}_\mathbb{C}^1 \rightarrow X$ algebraic morphism.

st. $f(0) = P_1$ and $f(1) = P_2$ (Stronger than path connected, 道路连通.)

(actually, $RC \Rightarrow \pi_1^{et}(X) = 0$. simply connected)

(2) This conjecture is of the style: geometry determines arithmetic.

(3) Conj \Rightarrow inverse of Galois problem.:

$\forall G$ finite gp, $\exists K/\mathbb{Q}$ finite Galois extension st:
 $Gal(K/\mathbb{Q}) = G.$

[it suffices to prove the conj. for $X =$ smooth compactification of SL_n/G]

(4) for non-rationally connected varieties,

\exists counter examples by Skorobogatov 90's
Poonen 2000's

(5) further obstructions?

Summary

